## Improving FIR Filter Coefficient Precision

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## Improving FIR Filter Coefficient Precision

- There is a method for increasing the precision of fixed-point coefficients used in linear-phase finite impulse response (FIR) filters,
- to achieve improved filter performance,
- without increasing either the number of coefficients or coefficient bitwidths.
- Thinking about this, such a process does not seem possible.
- But to see how, let's first review the behavior of a FIR filter.
- Consider an FIR filter's coefficients (impulse response) shown in Figure 1(a).
- Such a filter can be implemented as shown in Figure 1(b).


Figure 1

- Quantizing those coefficients in, an 8-bit two's complement format,
- yields the decimal integer and binary values in Figure 2(b).
- The beginning, and ending, coefficients are small in amplitude.
- Many high-order bits of the low-amplitude coefficients, the red-font underscored bits, are the same as the sign bit.
- That's because of the fixed bitwidth quantization.

(a)

Floating-point coefficients

| $b_{0}$ | 0.01751 | 2 | 00000010 |
| :---: | :---: | :---: | :---: |
| $b_{1}$ | -0.05899 | -8 | $1 \underline{1111000}$ |
| $b_{2}$ | -0.26156 | -33 | $1 \underline{1011111}$ |
| $b_{3}$ | 0.37687 | 48 | $0 \underline{0} 110000$ |
| $b_{4}$ | 0.87968 | 113 | 01110001 |
| $b_{5}$ | 0.37687 | 48 | $0 \underline{0} 110000$ |
| $b_{6}$ | -0.26156 | -33 | $1 \underline{1011111}$ |
| $b_{7}$ | -0.05899 | -8 | $1 \underline{1111000}$ |
| $b_{8}$ | 0.01751 | 2 | $\underline{000000} 10$ |

(b)

Figure 2

- Those underscored bits are "wasted" bits.
- They have no effect (no weight) on the calculation of filter output $y(n)$.

|  |  | sign bit |  |
| :---: | :---: | :---: | :---: |
| $b_{0}$ | 0.01751 | 2 | 00000010 |
| $b_{1}$ | -0.05899 | -8 | 11111000 |
| $b_{2}$ | -0.26156 | -33 | 11011111 |
| $b_{3}$ | 0.37687 | 48 | 00110000 |
| $b_{4}$ | 0.87968 | 113 | 01110001 |
| $b_{5}$ | 0.37687 | 48 | 00110000 |
| $b_{6}$ | -0.26156 | -33 | 11011111 |
| $b_{7}$ | -0.05899 | -8 | 11111000 |
| $b_{8}$ | 0.01751 | 2 | 00000010 |

- So the idea here is to replace those "wasted" bits with more significant bits,
- to give us improved numerical precision for the low-amplitude beginning and ending coefficients.
- OK, let's look at an example,
- of what's called a "serial" implementation of this whole idea.


## "Serial" Method

- Assume we quantize the maximum-amplitude coefficient, $\boldsymbol{b}_{4}$, to eight bits.
- Next, we quantize the lower-amplitude coefficients to larger bitwidths than the max-amplitude coefficient $\boldsymbol{b}_{\mathbf{4}}$.

| 8-bit quantization |  |  |  | Implied binary point |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $b_{0}$ | 0.01751 | (2) |  | 00000010.010 | 11-bit | (2.25) |
| $b_{1}$ | -0.05899 | (-8) |  | 11111000.10 | 10-bit | (-7.5) |
| $b_{2}$ | -0.26156 | (-33) |  | 11011110.1 | 9-bit | (-33.5) |
| $b_{3}$ | 0.37687 | (48) |  | 00110000.0 | 9-bit | (48) |
| $b_{4}$ | 0.87968 | (113) |  | 01110001. | 8-bit | (113) |
| $b_{5}$ | 0.37687 | (48) | bitwidths | 00110000.0 | 9-bit | (48) |
| $b_{6}$ | -0.26156 | (-33) | bitwidths | 11011110.1 | 9-bit | (-33.5) |
| $b_{7}$ | -0.05899 | (-8) |  | 11111000.10 | 10-bit | (-7.5) |
| $b_{8}$ | 0.01751 | (2) |  | 00000010.010 | 11-bit | (2.25) |

Figure 3

- We'll discuss how to choose the coefficients' variable bitwidths in a moment.
- Next we delete the appropriate "wasted" (red-underscored) bits,
- to arrive at our final 8-bit coefficients.

|  |  | Final coefficients |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 000000010010 |  | 00010010 | (18) | Flag $=1$ |
| 1111100010 |  | 11100010 | (-30) | Flag $=1$ |
| 110111101 |  | 10111101 | (-67) | Flag $=0$ |
| $0 \underline{01100000}$ |  | 01100000 | (96) | Flag $=1$ |
| 01110001 | (underscored) | - 01110001 | (113) | Flag $=0$ |
| 001100000 | bits | 01100000 | (96) | Flag $=1$ |
| 110111101 |  | 10111101 | (-67) | Flag $=0$ |
| 1111100010 |  | 11100010 | (-30) | Flag $=1$ |
| $0 \underline{0000010010}$ |  | 00010010 | (18) | Flag $=1$ |

Figure 4

- Appended to each coefficient is a flag bit,
- indicating whether that coefficient used one more quantization bit than the previous (next larger) coefficient.
- The question now is, "How do we use those "oddball" coefficients in a filter?"
- Figure 5 shows us the answer.
- This implementation is called "serial" because there is only one multiplier.
"Serial" filter implementation



## Figure 5

- For an $N$-tap FIR filter,
- for odd $N,(N+1) / 2$ coefficients are stored in the coefficient ROM (read-only memory).
- for even $N, N / 2$ coefficients are stored in the coefficient ROM.
- When a new $x(n)$ input sample arrives, we:
- Set the accumulator to zero.
- Multiply the sum of the appropriate data registers by the $\boldsymbol{b}_{4}$ coefficient.
- Add that product to the accumulator.
- Next we multiply the sum of the appropriate data registers by the $b_{3}$ coefficient.
-- If the flag bit of the $\boldsymbol{b}_{\mathbf{3}}$ coefficient is one, we left-shift the current accumulator value and add the current multiplier's output to the shifted accumulator value.
-- If the current coefficient's flag bit is zero the accumulator word is not shifted prior to an accumulation.
- Continue these multiplications, possible left shifts, and accumulations for the remaining $b_{2}, b_{1}$, and $b_{0}$ coefficients.
- So, when a new $x(n)$ input sample arrives, we perform a series of multiplications and accumulations (using multiple clock cycles),
- always starting with the largest coefficient $\left(b_{4}\right)$,
- to produce a single $y(n)$ filter output sample.
- To maintain our original FIR filter's gain,
- after the final accumulation we truncate the final accumulator value by discarding its least significant $M$ bits,
- where $M$ is the total number of flag bits in the ROM memory.
- Let's look at this "serial" method in action.


## "Serial" Example

- Implement a 29-tap lowpass FIR filter,
- whose cutoff frequency is $0.167 f_{s}$ and whose stopband begins at $0.292 f_{s}$.
(a)

(b)

Figure 6
- Relative to a traditional fixed-point implementation (dotted curve), the "serial method" (dashed curve) provides:
- Improved stopband attenuation,
- Reduced transition region width,
- Improved passband ripple performance.

- All of these improvements occur:
- without increasing the bitwidths of our filter's coefficients
- without increasing the number of coefficients.
- Regarding this "serial method", American actor Robert De Niro would say:
- I like it.
- I like it.
- What did I tell you?
- WHAT DID I TELL YOU?
- I like it!

- As it turns out, we can do even better than the "serial method".


## "Parallel" Method

- In the serial method, adjacent filter coefficients were quantized to a precision differing by no more that one bit.
- That's because we used "flag bits".
- In the parallel method, adjacent coefficients can be quantized to a precision differing by more than one bit.
- Figure 7 shows an example of our parallel method's coefficient quantization process.
- Again, assume we quantize the maximum-amplitude coefficient, $b_{4}$, to eight bits.
- Next, we quantize the lower-amplitude coefficients to larger bitwidths than the max-amplitude coefficient $\boldsymbol{b}_{\mathbf{4}}$.

| 8-bit quantization |  |  |  | Implied binary point |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $b_{0}$ | 0.01751 | (2) |  | 00000010.01000 | 13-bit | (2.25) |
| $b_{1}$ | -0.05899 | (-8) |  | 11111000.0111 | 12-bit | (-7.5625) |
| $b_{2}$ | -0.26156 | (-33) |  | 11011110.1 | 9-bit | (-33.5) |
| $b_{3}$ | 0.37687 | (48) |  | 00110000.0 | 9-bit | (48) |
| $b_{4}$ | 0.87968 | (113) | various | 01110001. | 8-bit | (113) |
| $b_{5}$ | 0.37687 | (48) | batwidths | 00110000.0 | 9-bit | (48) |
| $b_{6}$ | -0.26156 | (-33) | bitwidths | 11011110.1 | 9-bit | (-33.5) |
| $b_{7}$ | -0.05899 | (-8) |  | 11111000.0111 | 12-bit | (-7.5625) |
| $b_{8}$ | 0.01751 | (2) |  | 00000010.01000 | 13-bit | (2.25) |

## Figure 7

- Notice that $b_{2}$ is quantized to 9 bits, and
- $b_{1}$ is quantized to $\mathbf{1 2}$ bits.
- We'll discuss how to choose the coefficients' variable bitwidths in a moment.
- As before, we then delete the appropriate "wasted" (red-underscored) bits, - to arrive at our final 8 -bit coefficients.

|  |  | Final coefficients |  |
| :---: | :---: | :---: | :---: |
| 0000001001000 |  | 01001000 | (72) |
| 111110000111 |  | 10000111 | (-121) |
| 110111101 |  | 10111101 | (-67) |
| 0 O 1100000 | Delete wasted | 01100000 | (96) |
| 01110001 | (underscored) | > 01110001 | (113) |
| 001100000 | bits | 01100000 | (96) |
| 110111101 |  | 10111101 | (-67) |
| 111110000111 |  | 10000111 | (-121) |
| $0 \underline{000001001000}$ |  | 01001000 | (72) |

Figure 8

- Figure 9 shows the implementation of the "parallel" method.
- This implementation is called "parallel" because there are multiple multipliers.
- To keep our drawings simple, assume we're building a 5-tap filter.
- $b_{2}$ is the maximum-amplitude coefficient.


Figure 9

- When a new $x(n)$ input sample arrives, we:
- Set the accumulator to zero.
- Multiply the sums of the appropriate data registers by the corresponding coefficients.
-- All multiplications occur in one clock cycle (i.e., in parallel).
- The multiple products are added to the accumulator as shown in Figure 10.


Figure 10

- For example, if there were four wasted bits deleted from the high-precision $b_{1}$ coefficient,
-- then the $V_{k}$ product is shifted to the right by four bits, relative to the $\boldsymbol{W}_{k}$ product bits, before being added to the accumulator word.
- If there were seven wasted bits deleted from the high-precision $\boldsymbol{b}_{0}$ coefficient,
-- then the $\boldsymbol{U}_{k}$ product is shifted to the right by seven bits, relative to the $\boldsymbol{W}_{k}$ product bits, before being added to the accumulator word.
- It's the data routing that accounts for the deleted "wasted" bits in Figure 8!
- Let's look at this "parallel" method in action.


## "Parallel" Example

- Implementing the same 29-tap lowpass filter as in the "serial" method example yields the performance curves in Figure 11.


Figure 11

- Relative to the "serial method" (dashed curve) implementation, the "parallel" method (solid curve) provides:
- even further-improved stopband attenuation.
- Again, without increasing either the bitwidths of our filter's coefficients, or the number of coefficients.

- Siskel and Ebert would give this parallel method "Two Thumbs Up."


## Choosing the Number of Bits in Variable Bitwidth Coefficients

- There are algorithms for determining the number of bits in the variable bitwidth coefficients.
- One algorithm for the "serial" method coeffs. in Figure 3,
- and another algorithm for the "parallel" method coeffs. in Figure 7.
- Those algorithms are a bit too intricate (too grueling) to cover in a Conference presentation such as this.
- Those algorithms will be published in the "DSP Tips \& Tricks" column,
- in the July 2010 issue of the IEEE Signal Processing Magazine.
- If you want to learn those algorithms before July, send me an E-mail,
- at: [R.Lyons@ieee.org](mailto:R.Lyons@ieee.org).
- Please be aware that the Copyrights to the figures in this presentation are, this month, being transferred to the IEEE.
- This entire filter coefficient-enhancement idea is not mine.
- This is the idea of Zhi Shen.
- Ph.D degree student with the Department of Electronics and Information Engineering, Huazhong Univ.Sci. \& Tech., Wuhan, P.R. China.

- As far as I know, Mr. Shen has implemented these improved-precision coefficient methods,
- on an Altera FPGA.

